

# Fault Tolerant Quantum Error Correction

Master's Thesis Midterm  
Presentation

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# Overview

- Introduction to Quantum Error Correction
- Fault Tolerance and Detector Error Models
- Research Gap
- Simulation Methodology



# Quantum Computing

- Simulating quantum systems on classical hardware is exponentially complex  
→ Can't we use quantum hardware to simulate quantum systems? [Fey82]
- Some problems that are “hard” to solve on classical computers we can “easily” solve on quantum computers [Pre18]

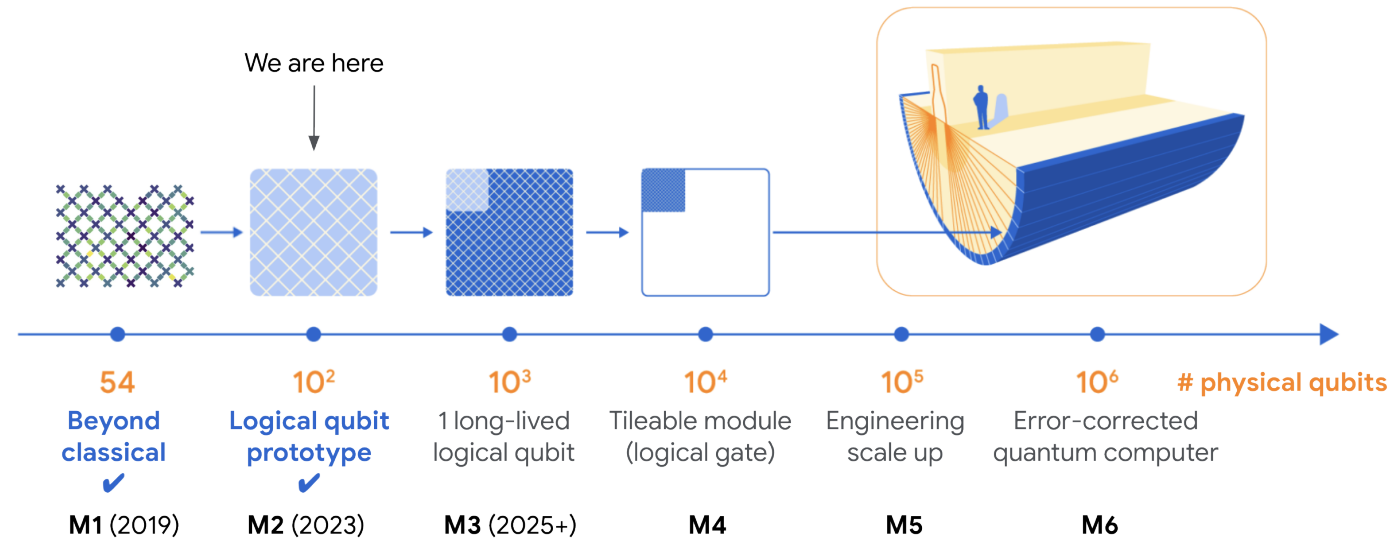


Figure: Google Quantum AI's quantum computing roadmap [Goo].

[Fey82] R. P. Feynman, “Simulating physics with computers,” en, *International Journal of Theoretical Physics*, vol. 21, no. 6, pp. 467–488, Jun. 1982,

[Pre18] J. Preskill, “Quantum Computing in the NISQ era and beyond,” en-GB, *Quantum*, vol. 2, p. 79, Aug. 2018,

[Goo] Google Quantum AI, *Quantum Computing Roadmap*, en. Accessed: Jan. 28, 2026.

# The Need for Quantum Error Correction



- Errors during quantum computation are inevitable because quantum systems are fragile
- We want to interact with the quantum state but not disturb it
- We employ more physical qubits to introduce redundancy and use the resulting *physical* state to represent the *logical* state [Rof19]
- Typical scales
  - IBM recently introduced a scheme encoding 12 logical qubits in 288 physical ones [BCG<sup>+</sup>24]
  - The physical error rate is typically assumed to be  $10^{-3}$  for simulations (e.g., [BCG<sup>+</sup>24])
  - Decoding has to happen with ultra-low latency to avoid the backlog problem (about 1 us per data extraction round) [CSB<sup>+</sup>24]

[Rof19] J. Roffe, “Quantum error correction: An introductory guide,” *Contemporary Physics*, vol. 60, no. 3, pp. 226–245, Jul. 2019,

[BCG<sup>+</sup>24] S. Bravyi et al., “High-threshold and low-overhead fault-tolerant quantum memory,” *Nature*, vol. 627, no. 8005, pp. 778–782, Mar. 2024,

[CSB<sup>+</sup>24] L. Caune et al., *Demonstrating real-time and low-latency quantum error correction with superconducting qubits*, Oct. 2024. Accessed: Jan. 28, 2026.

# Peculiarities of the Quantum Setting



- Quantum error correction (QEC) is actually able to protect the actual quantum state
- Similar to bits and gates, quantum systems are built on top of qubits and quantum gates
- We have to consider phase flip errors in addition to bit flip errors [Rof19]

$$\begin{aligned} |0\rangle &\rightarrow |1\rangle \\ |1\rangle &\rightarrow |0\rangle \end{aligned}$$

(a) Bit flip (X) error

$$\begin{aligned} |0\rangle &\rightarrow |0\rangle \\ |1\rangle &\rightarrow -|1\rangle \end{aligned}$$

(b) Phase flip (Z) error

$$\begin{aligned} |0\rangle &\rightarrow |1\rangle \\ |1\rangle &\rightarrow -j|0\rangle \end{aligned}$$

(c) Y error: Combination of X and Z

- Measuring the qubits directly destroys superpositions and entanglement  
→ We generally only work with the syndrome, which we can measure [NC10]
- Sometimes superposition permits multiple equivalent solutions to the decoding problem (*quantum degeneracy*) [RWBC20]

[NC10] M. A. Nielsen and I. L. Chuang, *Quantum Computation and Quantum Information: 10th Anniversary Edition*, en. Cambridge: Cambridge University Press, Dec. 2010, ISBN: 978-0-511-97666-7.

[Rof19] J. Roffe, “Quantum error correction: An introductory guide,” *Contemporary Physics*, vol. 60, no. 3, pp. 226–245, Jul. 2019,

[RWBC20] J. Roffe, D. R. White, S. Burton, and E. Campbell, “Decoding across the quantum low-density parity-check code landscape,” *Physical Review Research*, vol. 2, no. 4, p. 043423, Dec. 2020.

# Stabilizer and Calderbank Shor Steane Codes



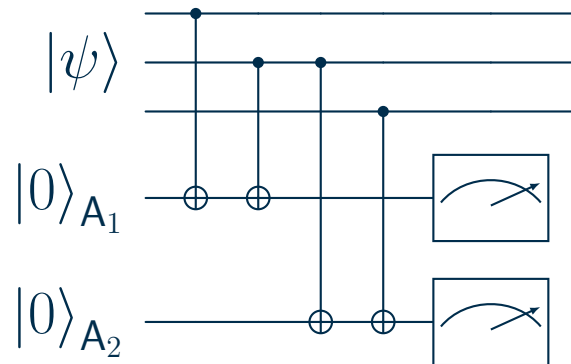
- Stabilizer codes [NC10]
  - The code space can implicitly be defined using *stabilizer generators*
  - We can represent them using parity check matrices
  - Quantum analog of linear codes
- Calderbank Shor Steane (CSS) codes [NC10]
  - Subset of stabilizer codes
  - Can correct X and Z errors independently
  - Described using two separate parity check matrices  $H_X$  and  $H_Z$
  - Can be constructed from two binary linear codes  $\mathcal{C}_1 [n, k_1]$  and  $\mathcal{C}_2 [n, k_2]$  with  $\mathcal{C}_2 \subset \mathcal{C}_1$

[NC10] M. A. Nielsen and I. L. Chuang, *Quantum Computation and Quantum Information: 10th Anniversary Edition*, en. Cambridge: Cambridge University Press, Dec. 2010, ISBN: 978-0-511-97666-7.

# Syndrome Extraction Circuits

- We entangle the state with *ancilla qubits* to perform syndrome measurements [NC10]
- Example: The 3-qubit repetition code<sup>1</sup>

$$\mathbf{H} = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}$$



[NC10] M. A. Nielsen and I. L. Chuang, *Quantum Computation and Quantum Information: 10th Anniversary Edition*, en. Cambridge: Cambridge University Press, Dec. 2010, ISBN: 978-0-511-97666-7.

<sup>1</sup>Note that, for simplicity, this chosen example is a code that is only able to correct X errors (bit flips)

# Overview



- Introduction to Quantum Error Correction
- Fault Tolerance and Detector Error Models
  - Fault Tolerance
  - Detector Error Models
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# Fault Tolerance

- The quantum gates we use for syndrome extraction are faulty themselves  
→ We need *fault-tolerant* QEC
- A QEC procedure is said to be fault tolerant if, in addition to correcting *input errors*, the spread of *internal errors* is sufficiently limited [DTTBE25]

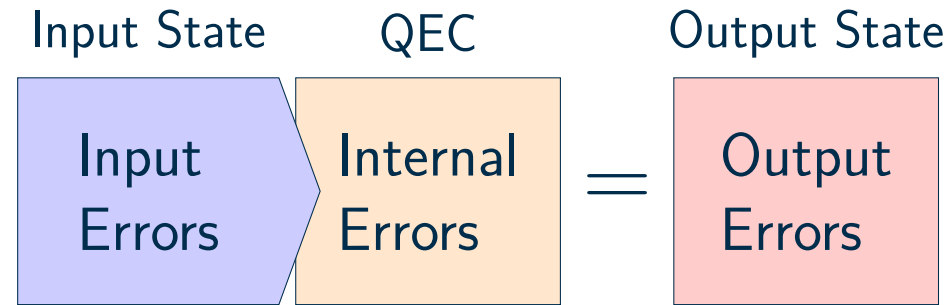


Figure: Overview of the flow of errors in a QEC system. Adapted from [DTTBE25].

- We have to modify the syndrome extraction circuitry to be fault tolerant (e.g., by using specially prepared multi-qubit states for each ancilla [Sho97])
- We generally perform multiple rounds of syndrome extraction

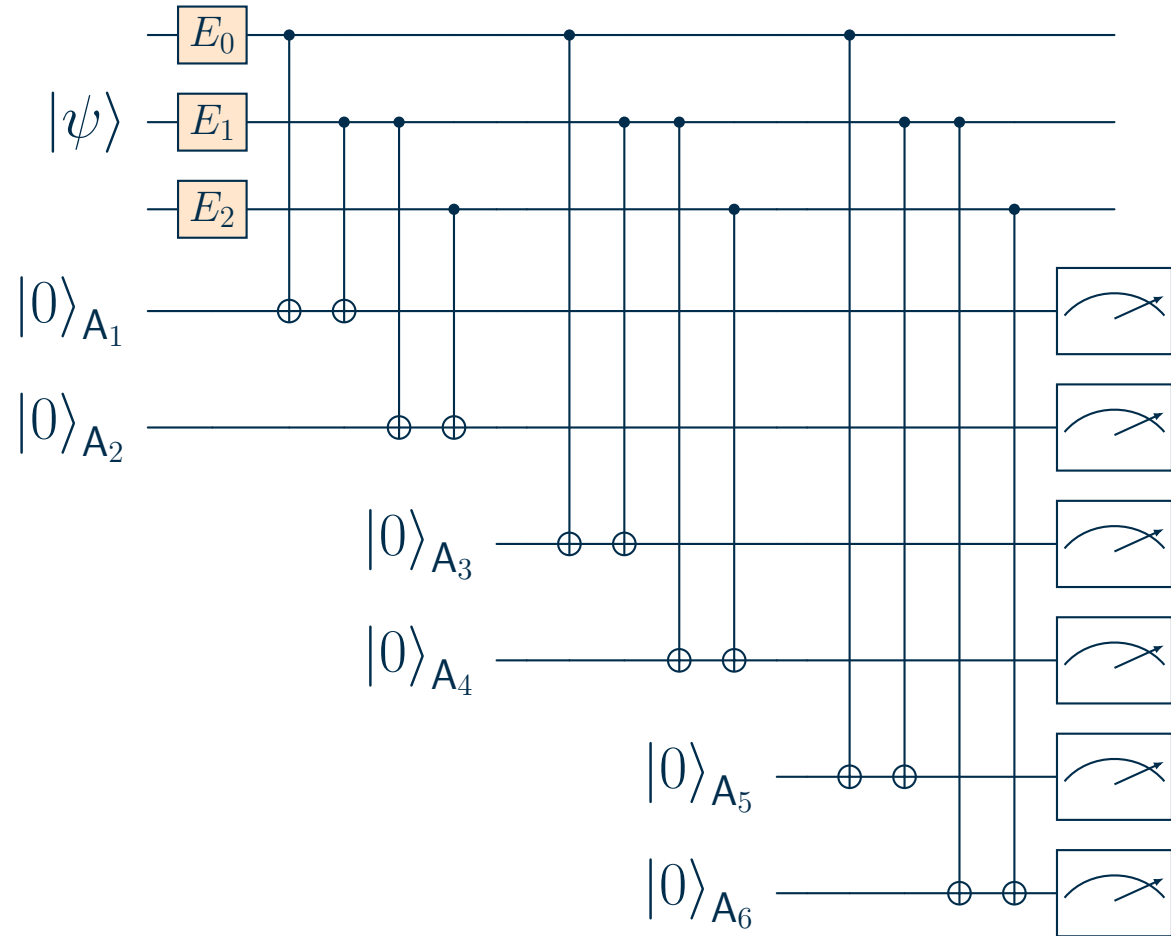
[Sho97] P. W. Shor, “Fault-tolerant quantum computation,” Mar. 1997,

[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

# The Measurement Syndrome Matrix I

- Each column of the *measurement syndrome matrix*  $\Omega$  corresponds to a measurement pattern an error produces [DTTBE25]
- Example: 3-qubit repetition code (Only bit flips on data qubits)

$$\Omega = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}$$

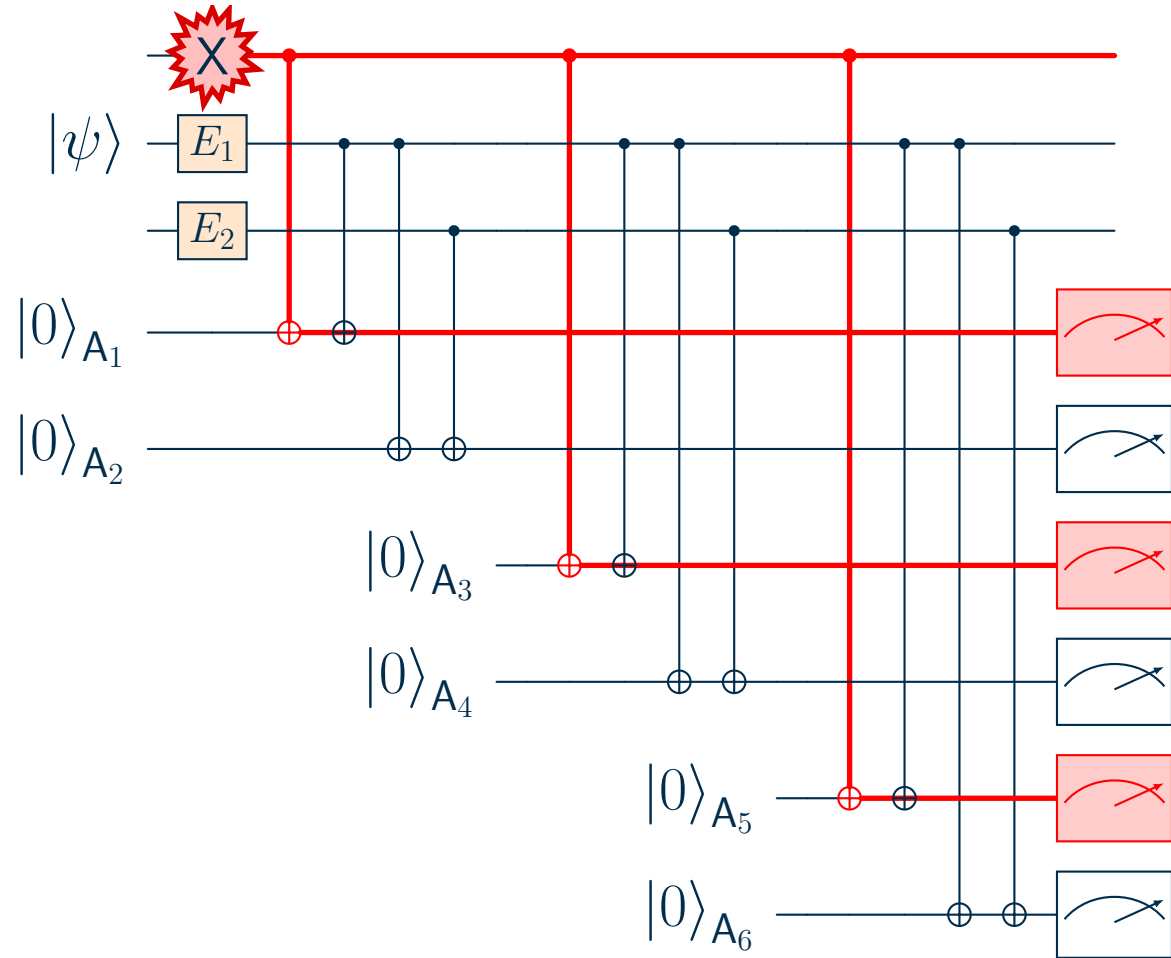


[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, "Designing fault-tolerant circuits using detector error models," Oct. 2025,

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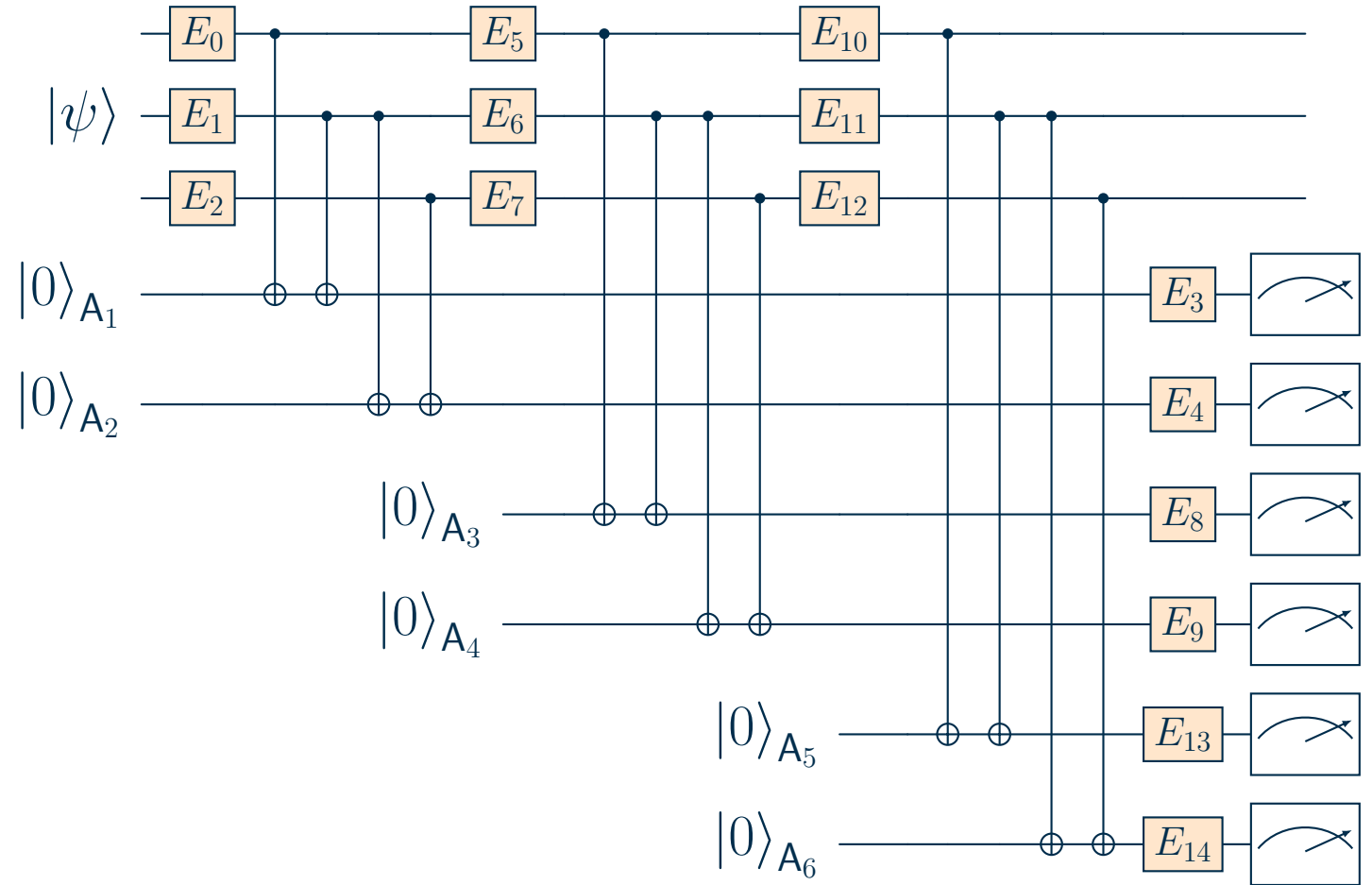


[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

# The Measurement Syndromemani Matrix II

- Each column of the *measurement syndrome matrix*  $\Omega$  corresponds to a measurement pattern an error produces [DTTBE25]
- Example: 3-qubit repetition code (Phenomenological noise [DTTBE25])

$$\Omega = \begin{pmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}$$

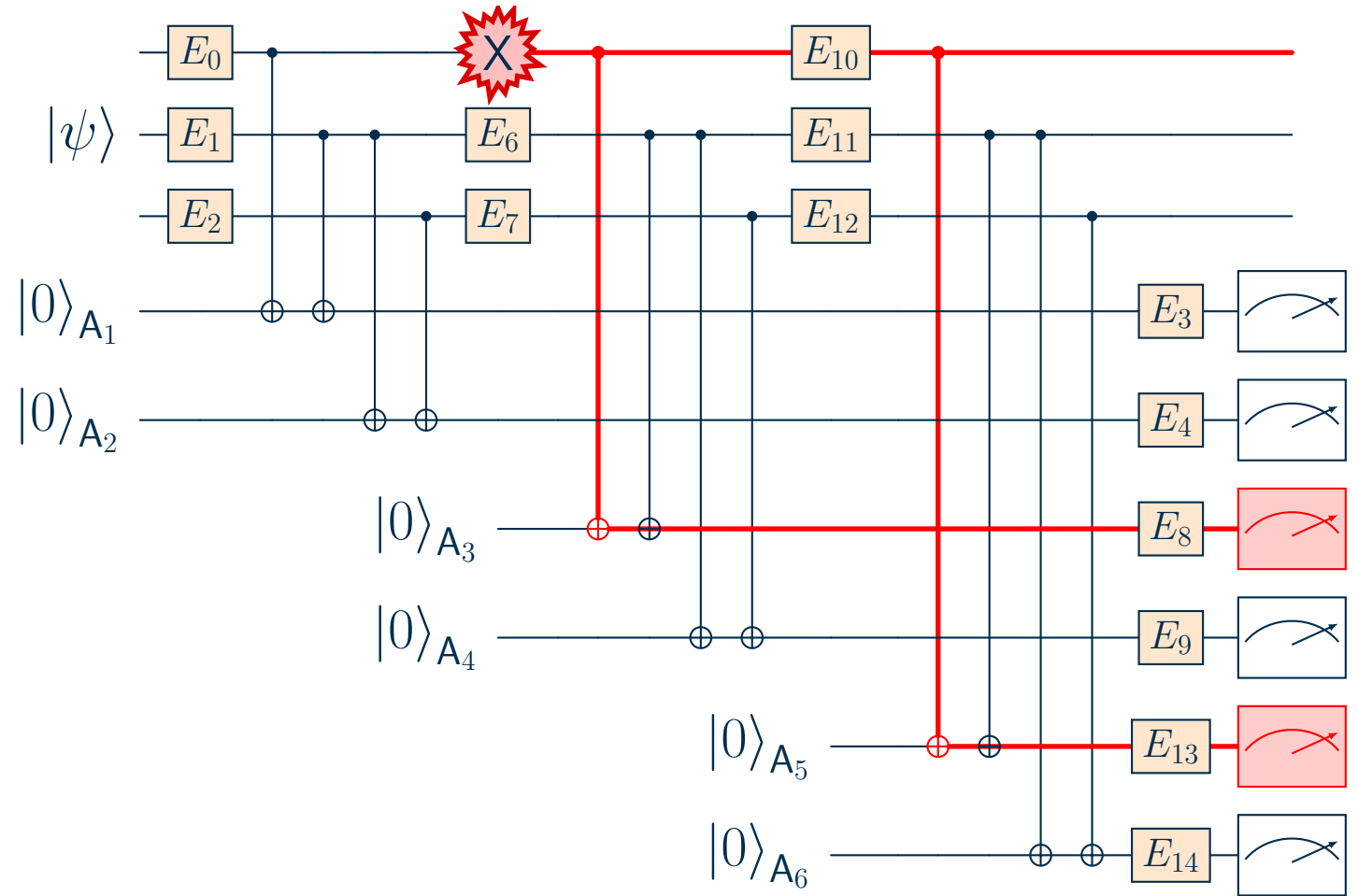


[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

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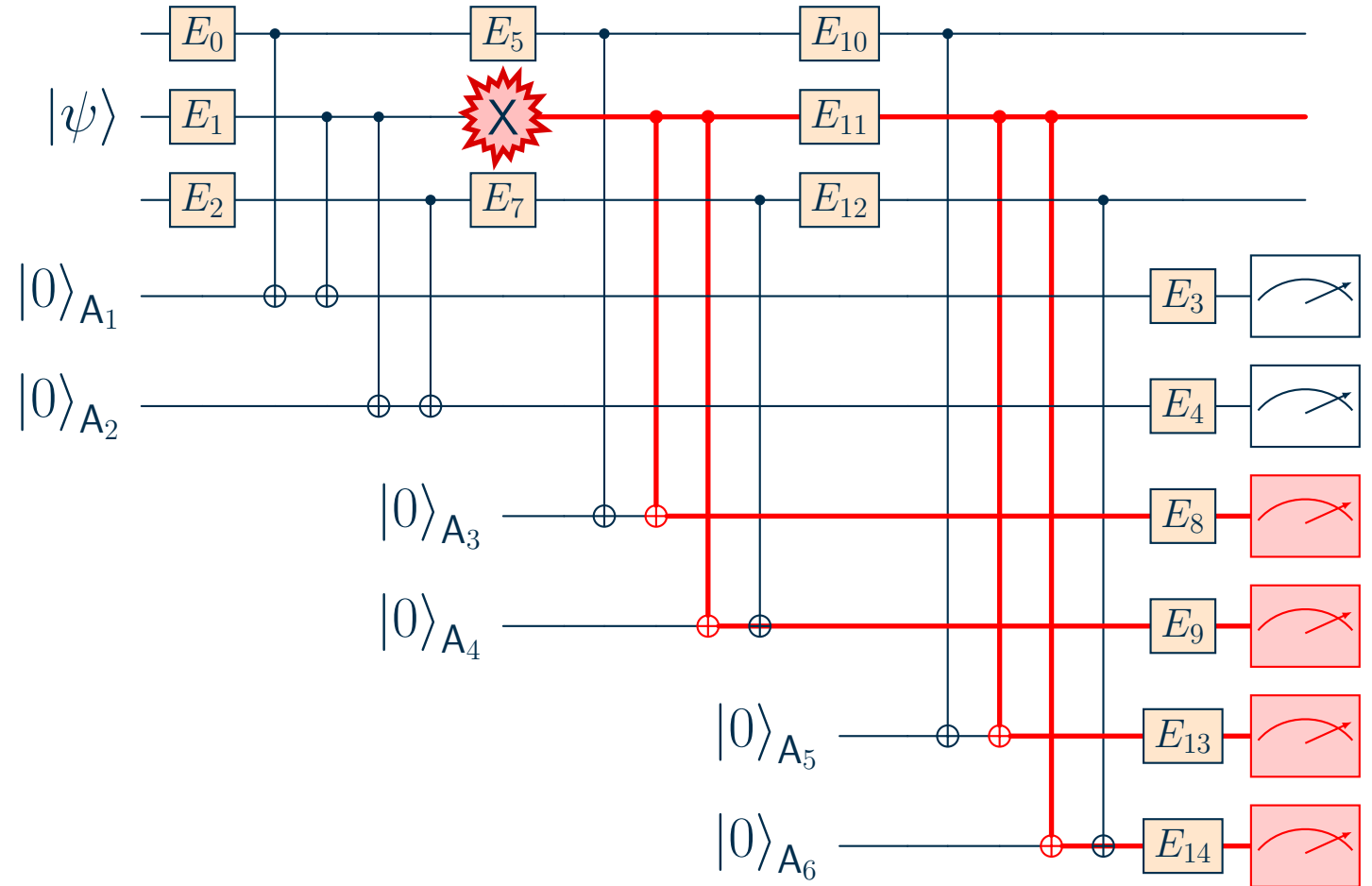


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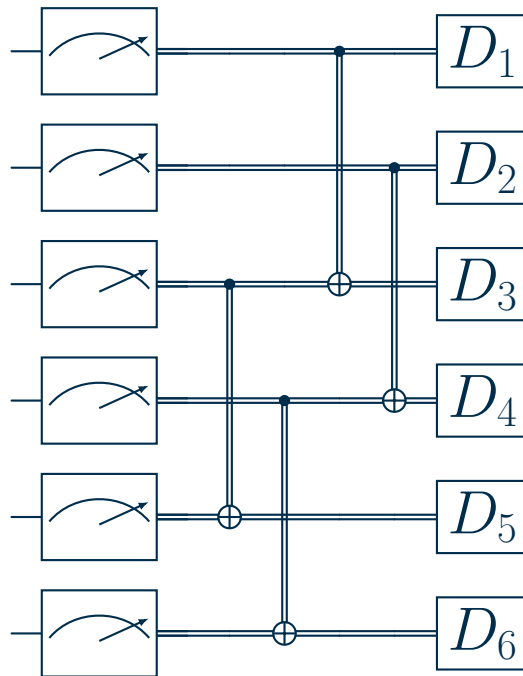
$$\Omega = \begin{pmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}$$



[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

# The Detector Error Matrix I

- A detector is a parity constraint on a set of measurement outcomes [DTTBE25]
- Each column of the *detector error matrix*  $H$  corresponds to a detector pattern an error produces
- We can mitigate the propagation of errors into subsequent rounds by XORing the measurements, i.e., defining detectors appropriately; this is equivalent to row additions

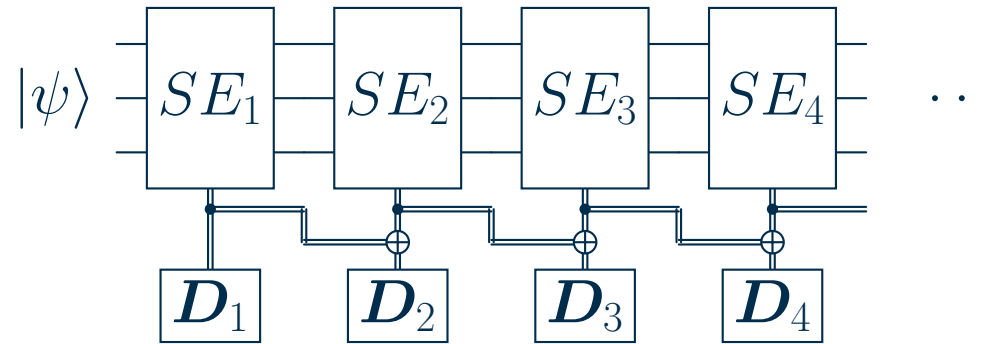


$$H = \begin{pmatrix} 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 \end{pmatrix}$$

[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

# The Detector Error Matrix II

- Visualization of general process



- E.g., for bivariate bicycle (BB) codes, the resulting detector error matrix under circuit-level noise has the form [\[GCR24\]](#)

$$H = \begin{pmatrix} H_0 & H_1 & 0 & 0 & 0 & 0 & \dots \\ 0 & H_2 & H_0 & H_1 & 0 & 0 & \\ 0 & 0 & 0 & H_2 & H_0 & H_1 & \\ 0 & 0 & 0 & 0 & 0 & H_2 & \\ \vdots & & & & & & \ddots \end{pmatrix}$$

[GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024. Accessed: Nov. 24, 2025.

# Noise Model Types

- The noise model assigns a likelihood to the occurrence of each error
- The *depolarizing channel* considers [NC10]
  - X, Y or Z errors on the data qubits
- *Phenomenological noise* considers [DTTBE25]
  - X errors on data qubits before each measurement round
  - X errors on measurement outcomes
- *Circuit-level noise* considers [DTTBE25]
  - X, Y or Z errors after state preparation
  - $n$ -qubit X, Y or Z errors after any  $n$ -qubit gate
  - X errors on measurement outcomes

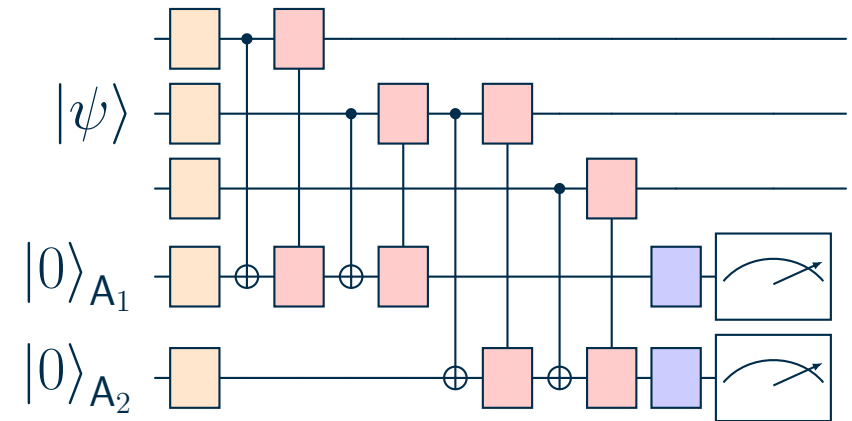


Figure: Circuit-level noise model for the 3-qubit repetition code (for X errors).

- [NC10] M. A. Nielsen and I. L. Chuang, *Quantum Computation and Quantum Information: 10th Anniversary Edition*, en. Cambridge: Cambridge University Press, Dec. 2010, ISBN: 978-0-511-97666-7.
- [DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

# Decoding using Detector Error Models



- A detector error model (DEM) combines a detector error matrix and a noise model
- When employing belief propagation (BP), the syndrome-based variant must be used [BBA<sup>+</sup>15]
- The likelihoods of different error locations can be used as priors for decoding
- Challenges
  - Repeated syndrome measurements come with increased decoding complexity [GCR24]
  - Degeneracy and short cycles lead to degraded performance of BP [BBA<sup>+</sup>15]

[BBA<sup>+</sup>15] Z. Babar et al., “Fifteen Years of Quantum LDPC Coding and Improved Decoding Strategies,” *IEEE Access*, vol. 3, pp. 2492–2519, 2015,  
[GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024.  
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- Research Gap
  - State of the Art
  - Proposed Research
- Simulation Methodology



# Addressing the Challenges



- Many window-based approaches exist to combat the decoding complexity
  - Parallel decoding of syndromes [SBB<sup>+</sup>23]
  - Using sliding windows [HP23] [GCR24]
- To deal with the degraded BP performance, it is usually modified or extended
  - Ordered statistics decoding (OSD) post-processing [RWBC20]
  - Guided decimation [GCR24]
  - Neural approaches [KL22] [MSLS25]
  - Ensemble decoding [KSWB25]

- [RWBC20] J. Roffe, D. R. White, S. Burton, and E. Campbell, “Decoding across the quantum low-density parity-check code landscape,” *Physical Review Research*, vol. 2, no. 4, p. 043423, Dec. 2020.
- [KL22] K.-Y. Kuo and C.-Y. Lai, “Exploiting degeneracy in belief propagation decoding of quantum codes,” en, *npj Quantum Information*, vol. 8, no. 1, p. 111, Sep. 2022,
- [HP23] S. Huang and S. Puri, “Improved Noisy Syndrome Decoding of Quantum LDPC Codes with Sliding Window,” Nov. 2023,
- [SBB<sup>+</sup>23] L. Skoric et al., “Parallel window decoding enables scalable fault tolerant quantum computation,” en, *Nature Communications*, vol. 14, no. 1, p. 7040, Nov. 2023,
- [GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024. Accessed: Nov. 24, 2025.
- [MSLS25] S. Miao, A. Schnerring, H. Li, and L. Schmalen, “Quaternary Neural Belief Propagation Decoding of Quantum LDPC Codes with Overcomplete Check Matrices,” *IEEE Access*, vol. 13, pp. 25 637–25 649, Feb. 2025,
- [KSWB25] S. Koutsoumpas, H. Sayginel, M. Webster, and D. E. Browne, “Automorphism Ensemble Decoding of Quantum LDPC Codes,” en, Mar. 2025

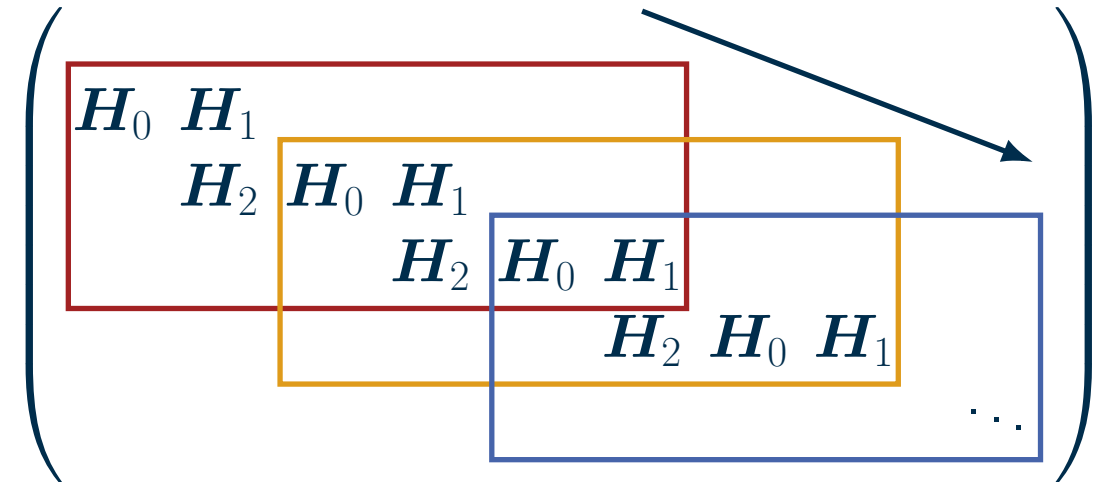
# Sliding-Window Decoding

- The sliding window approach taken in [GCR24] resembles a spatially coupled low density parity check (SC-LDPC) code decoder
- However, they don't pass soft information between windows (only update of syndrome based on hard decision in previous window)

$$\begin{pmatrix} H_0 & H_1 & 0 & 0 & 0 & 0 \\ 0 & H_2 & H_0 & H_1 & 0 & 0 \\ 0 & 0 & 0 & H_2 & H_0 & H_1 \end{pmatrix} \begin{pmatrix} \hat{e}_0 \\ \vdots \\ \hat{e}_5 \end{pmatrix} = \begin{pmatrix} s_1 \\ s_2 \\ s_3 \end{pmatrix}$$

$$s'_2 = s_2 + H_2 \hat{e}_1$$

(a) Equations for the decoding of the first window



(b) Visualization of sliding window procedure

- They try BP + OSD and a modification of BP with guided decimation

[GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024. Accessed: Nov. 24, 2025.

- Research gap
  - Existing literature into circuit-level noise fails to properly consider the SC-LDPC like structure of the detector error matrix
- Proposed Methodology
  - Adapt modified guided decimation decoder from [GCR24] to pass soft information
  - Investigate performance of different modifications of BP for "inner decoder" (e.g., quaternary neural BP [MSLS25])

[GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024. Accessed: Nov. 24, 2025.

[MSLS25] S. Miao, A. Schnerring, H. Li, and L. Schmalen, "Quaternary Neural Belief Propagation Decoding of Quantum LDPC Codes with Overcomplete Check Matrices," *IEEE Access*, vol. 13, pp. 25 637–25 649, Feb. 2025,

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- **Simulation Methodology**
  - The Code and Other Parameters
  - Proposed Methodology



# Memory and Stability experiments



- What is a memory experiment? Communications Engineering view (what are my inputs and output? What do I expect?)
- What is a stability experiment?

# Noise Models and Figures of Merit



- For circuit-level noise, often, all error probabilities are set to the same value for simulations [FSG09]
- There are other approaches (e.g., SDMB noise, SI noise) [DTTBE25]
- Footprint plots
- Other figure of merit (Look into ECCentric?)

[DTTBE25] P.-J. H. S. Derks, A. Townsend-Teague, A. G. Burchards, and J. Eisert, “Designing fault-tolerant circuits using detector error models,” Oct. 2025,

[FSG09] A. G. Fowler, A. M. Stephens, and P. Groszkowski, “High-threshold universal quantum computation on the surface code,” *Physical Review A*, vol. 80, no. 5, p. 052312, Nov. 2009,

# Proposed Simulation Methodology



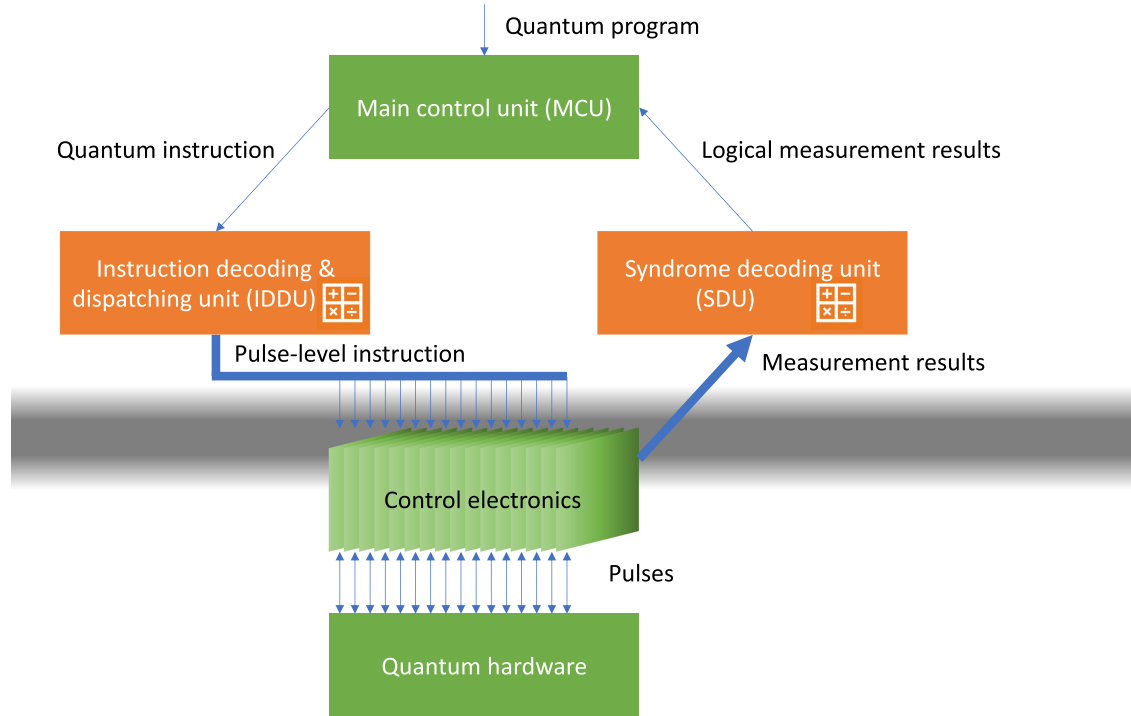
- Noise model
- Memory or stability experiment
- Figure of merit: Footprint plot
- Comparison with BB code also simulated by [GCR24]
- Comparison with surface code

[GCR24] A. Gong, S. Cammerer, and J. M. Renes, *Toward Low-latency Iterative Decoding of QLDPC Codes Under Circuit-Level Noise*, en, Mar. 2024.  
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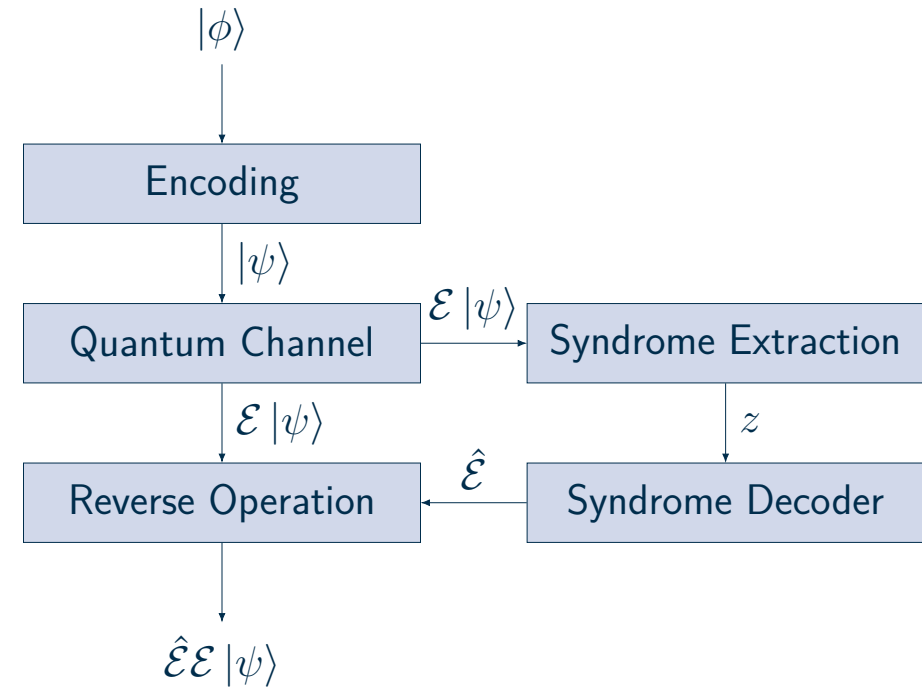
Thank you for your attention!  
Any questions?



# System Level Overview



(a) Schematic workflow of surface code quantum computation [ZZC<sup>+</sup>23].



(b) Block diagram of QEC using stabilizer codes [MSLS25].

[ZZC<sup>+</sup>23] F. Zhang et al., “A Classical Architecture for Digital Quantum Computers,” *ACM Transactions on Quantum Computing*, vol. 5, no. 1, 3:1–3:24, Dec. 2023,

[MSLS25] S. Miao, A. Schnerring, H. Li, and L. Schmalen, “Quaternary Neural Belief Propagation Decoding of Quantum LDPC Codes with Overcomplete Check Matrices,” *IEEE Access*, vol. 13, pp. 25 637–25 649, Feb. 2025,